EE382N (20): Computer Architecture - Parallelism and Locality Fall 2011

Lecture 17 – GPUs (II)

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Recap

Streaming model

- 1. Use many "slimmed down cores" to run in parallel
- 2. Pack cores full of ALUs (by sharing instruction stream across groups of fragments)
 - Option 1: Explicit SIMD vector instructions
 - Option 2: Implicit sharing managed by hardware
- 3. Avoid latency stalls by interleaving execution of many groups of fragments

- When one group stalls, work on another group Kayvon Fatahaliane 382N: Principles of Computer Architecture, Fall 2011 -- Lecture 17 (c) Mattan Erez Kayvon Fatahal

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Make the Compute Core The Focus of the 3 Architecture

- Photessessing
- Sterbuild the perbitient une adeusplet billip ralsessor computing



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Next-Gen GPU Architecture: Fermi



- 3 billion transistors
 - Over 2x the cores (512 total)
- ~2x the memory bandwidth

L1 and L2 caches

- 8x the peak DP performance
 - ECC
 - C++

Announced Sept. 2009

Fermi Focus Areas

- Expand performance sweet spot of the GPU
 - Caching
 - Concurrent kernels
 - FP64
 - 512 cores
 - GDDR5 memory
- Bring more users, more applications to the GPU
 - C++
 - Visual Studio
 Integration
 - ECC



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Streaming Multiprocessor (SM)

- Objective optimize for GPU computing Instruct
 - New ISA
 - Revamp issue / control flow
 - New CUDA core architecture
- 16 SMs per Fermi chip
- 32 cores per SM (512 total)
- 64KB of configurable
 L1\$ / shared memory

	FP32	FP64	INT	SFU	LD/ST
Ops / clk	32	16	32	4	16



Instruction Cache				
Scheduler		Scheduler		
Dispatch		Dispatch		
Register File				
Core	Core	Core	Core	
Core	Core	Core	Core	
Core	Core	Core	Core	
Core	Core	Core	Core	
Core	Core	Core	Core	
Core	Core	Core	Core	
Core	Core	Core	Core	
Core	Core	Core	Core	
Load/Store Units x 16				
Special Func Units x 4				
Interconnect Network				
64K Configurable Cache/Shared Mem				
Uniform Cache				

SM Microarchitecture

- New IEEE 754-2008 arithmetic standard
- Fused Multiply-Add (FMA) for SP & DP
- New integer ALU optimized for 64-bit and extended precision ops





Instruction Cache					
Scheduler		Scheduler			
Dispatch		Dispatch			
Register File					
Core	Core	Core	Core		
Core	Core	Core	Core		
Core	Core	Core	Core		
Core	Core	Core	Core		
Core	Core	Core	Core		
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Core	Core	Core	Core		
Core	Core	Core	Core		
Load	Load/Store Units x 16				
Spec	Special Func Units x 4				
Interconnect Network					
64K Configurable Cache/Shared Mem					
Uniform Cache					

Memory Hierarchy

- True cache hierarchy + on-chip shared RAM
 - On-chip shared memory: good fit for regular memory access
 - dense linear algebra, image processing, ...
 - Caches: good fit for irregular or unpredictable memory access
 - ray tracing, sparse matrix multiply, physics ...
- Separate L1 Cache for each SM (16/48 KB)
 - Improves bandwidth and reduces latency
- Unified L2 Cache for all SMs (768 KB)
 - Fast, coherent data sharing across all cores in the GPU





Larger, Faster Memory Interface

- GDDR5 memory interface
 - 2x improvement in peak speed over GDDR3
- Up to 1 Terabyte of memory attached to GPU
 - Operate on large data sets





GigaThread[™] Hardware Thread Scheduler

 Hierarchically manages tens of thousands of simultaneously active threads

- 10x faster context switching on Fermi
- Overlapping kernel execution







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GigaThread Streaming Data Transfer Engine

- Dual DMA engines
- Simultaneous CPU→GPU and GPU→CPU data transfer



 Fully overlapped with CPU/GPU processing Kernel 0 SDT1 CPU GPU SDT0 Kernel 1 CPU SDT0 **GPU** SDT1 Kernel 2 CPU SDT0 **GPU** SDT1 Kernel 3 CPU SDT0 GPU SDT1

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Other Capabilities

- ECC protection for DRAM, L2, L1, RF
- Unified 40-bit address space for local, shared, global
- 5-20x faster atomics
- ISA extensions for C++ (e.g. virtual functions)



	G80	GT200	Fermi	13
Transistors	681 million	1.4 billion	3.0 billion	
CUDA Cores	128	240	512	
Double Precision Floating Point	-	30 FMA ops/clock	256 FMA ops/clock	
Single Precision Floating Point	128 MAD ops/clock	240 MAD ops/clock	512 FMA ops/clock	
Special Function Units (per SM)	2	2	4	
Warp schedulers (per SM)	1	1	2	
Shared Memory (per SM)	16 KB	16 KB	Configurable 48/16 KB	
L1 Cache (per SM)	-	-	Configurable 16/48 KB	
L2 Cache	-	-	768 KB	
ECC Memory Support	-	-	Yes	
Concurrent Kernels	-	-	Up to 16	
Load/Store Address Width	32-bit	32-bit	64-bit	
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NVIDIA Tesla GPUs Power 3 of Top 5 Supercomputers

7168 Tesla GPUs 2.5 PFLOPS 4650 Tesla GPUs 1.2 PFLOPS

4224 Tesla GPUs 1.194 **PFLOPS**

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8 more GPU accelerated machines in the November **Top500**





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GPU Supercomputers: More Power Efficient



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Compute Core



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GeForce-8 Series HW Overview



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CUDA Processor Terminology

- **SPA** Streaming Processor Array
 - Array of TPCs
 - 8 TPCs in GeForce8800
- TPC Texture Processor Cluster
 - Cluster of 2 SMs + 1 TEX
 - TEX is a texture processing unit
- **SM** Streaming Multiprocessor
 - Array of 8 SPs
 - Multi-threaded processor core
 - Fundamental processing unit for a thread block
- **SP** Streaming Processor
 - Scalar ALU for a single thread

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With 1K of registers

Streaming Multiprocessor (SM)

- Streaming Multiprocessor (SM)
 - 8 Streaming Processors (SP)
 - 2 Super Function Units (SFU)
- Multi-threaded instruction dispatch
 - Vectors of 32 threads (warps)
 - Up to 16 warps per thread block
 - HW masking of inactive threads in a warp
 - Threads cover latency of texture/memory loads
- 20+ GFLOPS
- 16 KB shared memory
- 32 KB in registers
- DRAM texture and memory access

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Thread Life Cycle in HW

Kernel is launched on the SPA Host Device Kernels known as grids of thread blocks Grid 1 Thread Blocks are serially distributed Kernel **Block** 1 to all the SM's (0, 0)Potentially >1 Thread Block per SM At least 96 threads per block Each SM launches Warps of Thread 2 levels of parallelism Kernel 2 SM schedules and executes Warps that are ready to run **Block (1, 1)** As Warps and Thread Blocks Thread Thread complete, resources are freed (0, 0)(1, 0)SPA can distribute more Thread Blocks Thread Thread (1, 1)(0, 1)Thread Thread

Block Block Block (1, 1)(0, 1)(2, 1) Grid 2 Thread Thread Thread (2, 0)(3, 0)(4, 0)Thread Thread Thread (2.1)(3.1)(4, 1)Thread Thread Thread (0, 2)(1, 2)(2, 2)(3, 2)(4, 2)

Block

(1, 0)

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Block

(2, 0)

SM Executes Blocks



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Blocks

- Threads are assigned to SMs in Block granularity
 - Up to 8 Blocks to each SM as resource allows
 - SM in G80 can take up to 768 threads
 - Could be 256 (threads/block) * 3 blocks
 - Or 128 (threads/block) * 6 blocks, etc.
 - Threads run concurrently
 - SM assigns/maintains thread IDs
 - SM manages/schedules thread execution

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Thread Scheduling/Execution

- Each Thread Block is divided into 32-thread Warps
 - This is an implementation decision
- Warps are scheduling units in SM
- If 3 blocks are assigned to an SM and each Block has 256 threads, how many Warps are there in an SM?
 - Each Block is divided into 256/32 = 8 Warps
 - There are 8 * 3 = 24 Warps
 - At any point in time, only one of the 24 Warps will be selected for instruction fetch and execution.



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SM Warp Scheduling

- SM hardware implements zerooverhead Warp scheduling
 - Warps whose next instruction has its operands ready for consumption are eligible for execution
 - All threads in a Warp execute the same instruction when selected
 - Scoreboard scheduler
- 4 clock cycles needed to dispatch the same instruction for all threads in a Warp in G80
 - If one global memory access is needed for every 4 instructions
 - A minimal of 13 Warps are needed to fully tolerate 200-cycle memory latency



SM multithreaded Warp scheduler

time warp 8 instruction 11 warp 1 instruction 42 warp 3 instruction 95 warp 8 instruction 12 warp 8 instruction 12 warp 3 instruction 96

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SM Instruction Buffer – Warp Scheduling

- Fetch one warp instruction/cycle
 - from instruction L1 cache
 - into any instruction buffer slot
- Issue one "ready-to-go" warp instruction/cycle
 - from any warp instruction buffer slot
 - operand scoreboarding used to prevent hazards
- Issue selection based on round-robin/age of warp
- SM broadcasts the same instruction to 32 Threads of a Warp

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Scoreboarding

- All register operands of all instructions in the Instruction Buffer are scoreboarded
 - Status becomes ready after the needed values are deposited
 - prevents hazards
 - cleared instructions are eligible for issue
- Decoupled Memory/Processor pipelines
 - any thread can continue to issue instructions until scoreboarding prevents issue
 - allows Memory/Processor ops to proceed in shadow of Memory/Processor <u>ops</u>



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Granularity and Resource Considerations

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- For Matrix Multiplication, should I use 8X8, 16X16 or 32X32 tiles (1 thread per tile element)?
 - For 8X8, we have 64 threads per Block. Since each SM can take up to 768 threads, it can take up to 12 Blocks.
 However, each SM can only take up to 8 Blocks, only 512 threads will go into each SM!
 - For 16X16, we have 256 threads per Block. Since each SM can take up to 768 threads, it can take up to 3 Blocks and achieve full capacity unless other resource considerations overrule.

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- For 32X32, we have 1024 threads per Block. Not even one Kirk/NVIDIA and ei W. Hwu, 2007 1to an SM!

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Parallel Computing on a GPU

- NVIDIA GPU Computing Architecture
- Via a separate HW interface
- In laptops, desktops, workstations, servers
- 8-series GPUs deliver 50 to 200 GFLOPS on compiled parallel C applications
- GPU parallelism is doubling every year
- Programming model scales transparently
- Programmable in C with CUDA tools
- Multithreaded SPMD model uses application
 data parallelism and thread parallelism



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Next Lectures

- NVIDIA GeForce 8800 architecture
- CUDA programming model

